

MODULAR COLORINGS OF CORONA PRODUCT OF  $C_m$  WITH  $C_n$ 

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ABSTRACT. For  $\ell \geq 2$ , a modular  $\ell$ -coloring of a graph  $\mathcal{G}$  without isolated vertices is a coloring of the vertices of  $\mathcal{G}$  with the elements in  $\mathbb{Z}_\ell$  having the property that for every two adjacent vertices of  $\mathcal{G}$ , the sums of the colors of their neighbors are different in  $\mathbb{Z}_\ell$ . The minimum  $\ell$  for which  $\mathcal{G}$  has a modular  $\ell$ -coloring is the modular chromatic number of  $\mathcal{G}$ . In this paper, we determine the modular chromatic number of corona product of cycles.

## 1. INTRODUCTION

For graph-theoretical terminology and notation, we in general follow [1]. For a vertex  $v$  of a graph  $\mathcal{G}$ , let  $N_{\mathcal{G}}(v)$ , the *neighborhood of  $v$* , denote the set of vertices adjacent to  $v$  in  $\mathcal{G}$ . For a graph  $\mathcal{G}$  without isolated vertices, let  $c : V(\mathcal{G}) \rightarrow \mathbb{Z}_\ell$ ,  $\ell \geq 2$ , be a vertex coloring of  $\mathcal{G}$  where adjacent vertices may be colored the same. The *color sum*  $\mathcal{S}(v) = \sum_{u \in N_{\mathcal{G}}(v)} c(u)$  of a vertex  $v$  of  $\mathcal{G}$  is the sum of the colors of the vertices in  $N_{\mathcal{G}}(v)$ . The coloring  $c$  is called a *modular  $\ell$ -coloring* of  $\mathcal{G}$  if  $\mathcal{S}(x) \neq \mathcal{S}(y)$  in  $\mathbb{Z}_\ell$  for all pairs  $x, y$  of adjacent vertices in  $\mathcal{G}$ . The *modular chromatic number*  $Mc(\mathcal{G})$  of  $\mathcal{G}$  is the minimum  $\ell$  for which  $\mathcal{G}$  has a modular  $\ell$ -coloring. This concept was introduced by Zhang et. al. [2].

Okamoto, Salehi and Zhang proved, in [2], they proved that: every nontrivial connected graph  $\mathcal{G}$  has a modular  $\ell$ -coloring for some integer  $\ell \geq 2$  and  $Mc(\mathcal{G}) \geq \chi(\mathcal{G})$ , where  $\chi(\mathcal{G})$  denotes the chromatic number of  $\mathcal{G}$ ; for the cycle  $C_n$  of length  $n$ ,  $Mc(C_n)$  is 2 if  $n \equiv 0 \pmod{4}$  and it is 3 otherwise; every nontrivial

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tree has modular chromatic number 2 or 3; for the complete multipartite graph  $\mathcal{G}$ ,  $Mc(\mathcal{G}) = \chi(\mathcal{G})$ ; for the cartesian product  $\mathcal{G} = K_r \square K_2$ ,  $Mc(\mathcal{G})$  is  $r$  if  $r \equiv 2 \pmod{4}$  and it is  $r+1$  otherwise; for the wheel  $W_n = C_n \vee K_1$ ,  $n \geq 3$ ,  $Mc(W_n) = \chi(W_n)$ , where  $\vee$  denotes the join of two graphs; for  $n \geq 3$ ,  $Mc(C_n \vee K_2^c) = \chi(C_n \vee K_2^c)$ , where  $\mathcal{G}^c$  denotes the complement of  $\mathcal{G}$ ; and for  $n \geq 2$ ,  $Mc(P_n \vee K_2) = \chi(P_n \vee K_2)$ , where  $P_n$  denotes the path of length  $n-1$ ; and in [3] proved that: for  $m, n \geq 2$ ,  $Mc(P_m \square P_n) = 2$ .

Paramaguru and Sampthkumar proved, in [5], that:  $Mc(C_3 \square P_2) = 4$ ; except some special cases, for  $m \geq 3$  and  $n \geq 2$ ,  $Mc(C_m \square P_n) = \chi(C_m \square P_n)$ ; if  $m \equiv 2 \pmod{4}$  and  $n \equiv 1 \pmod{4}$ , then  $Mc(C_m \square P_n) \leq 3$ ; if  $n \equiv 1 \pmod{4}$ , then  $Mc(C_6 \square P_n) = 3$ . In [6], they proved that: if  $m \geq 4$  and  $n \geq 4$  are even integers and at least one of  $m, n$  is congruent to  $0 \pmod{4}$ , then  $Mc(C_m \square C_n) = \chi(C_m \square C_n)$ ; if  $n \geq 3$  is an integer, then  $Mc(C_3 \square C_n) = \chi(C_3 \square C_n)$ ; if at least one of  $m, n$  is congruent to  $1 \pmod{2}$ , except some special cases,  $m \geq 4$ ,  $n \geq 4$ , then  $Mc(C_m \square C_n) = \chi(C_m \square C_n)$ ; if  $n \equiv 2 \pmod{4}$ , and  $n \geq 6$ , then  $Mc(C_6 \square C_n) = 3$ , where  $\square$  denotes the Cartesian product of two graphs.

Nicholas and Sanma discussed in [4], that: the modular chromatic number of Fan, Helm graph, Friendship graph and gear graph.

The *corona* of two graphs  $\mathcal{G}$  and  $\mathcal{H}$  is the graph  $\mathcal{G} \circ \mathcal{H}$  formed from one copy of  $\mathcal{G}$  and  $|V(\mathcal{G})|$  copies of  $\mathcal{H}$ , where the  $i$ th vertex of  $\mathcal{G}$  is adjacent to every vertex in the  $i$ th copy of  $\mathcal{H}$ . Such type of graph products was introduced by Frucht and Harary in 1970.

## 2. CORONA OF $C_m$ WITH $C_n$

Define  $V(C_m) = \{u_1, u_2, u_3, \dots, u_m\}$ ;  $V(C_n) = \{v_1, v_2, v_3, \dots, v_n\}$ ;  $E(C_m) = \{u_1u_2, u_2u_3, u_3u_4, \dots, u_{m-1}u_m, u_mu_1\}$ ;  $E(C_n) = \{v_1v_2, v_2v_3, v_3v_4, \dots, v_{n-1}v_n, v_nv_1\}$ ;  $V(C_m \circ C_n) = V(C_m) \cup \{v_j^i : i \in \{1, 2, 3, \dots, m\} \text{ and } j \in \{1, 2, 3, \dots, n\}\}$ ;  $E(C_m \circ C_n) = E(C_m) \cup \{v_j^i v_{j+1}^i : i \in \{1, 2, 3, \dots, m\} \text{ and } j \in \{1, 2, 3, \dots, n-1\}\} \cup \{u_i v_j^i : i \in \{1, 2, 3, \dots, m\} \text{ and } j \in \{1, 2, 3, \dots, n\}\} \cup \{v_n^i v_1^i : i \in \{1, 2, 3, \dots, m\}\}$ .

**Theorem 2.1.** For  $m$  even and  $n$  even,  $m \geq 4$ ,  $n \geq 4$ ,  $Mc(C_m \circ C_n) = 3$ .

*Proof.* Let  $c : V(C_m \circ C_n) \rightarrow \mathbb{Z}_3$ .

**Case 1.**  $n \equiv 4 \pmod{6}$ .

Define  $c$  as follows:  $c(u_i) = 0$  if  $i$  is even;  $c(u_i) = 1$  if  $i$  is odd;  $c(v_j^i) = 0$  if

$i \in \{1, 2, 3, \dots, m\}$ ,  $j$  is even;  $c(v_j^i) = 1$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j$  is odd; then  $\mathcal{S}(u_i) = 1$  if  $i$  is even;  $\mathcal{S}(u_i) = 2$  if  $i$  is odd;  $\mathcal{S}(v_j^i) = 1$  if  $i, j$  odd;  $\mathcal{S}(v_j^i) = 2$  if  $i, j$  even;  $\mathcal{S}(v_j^i) = 0$  if  $i$  is odd,  $j$  is even;  $\mathcal{S}(v_j^i) = 0$  if  $i$  is even,  $j$  is odd.

**Case 2.**  $n \equiv 2 \pmod{6}$ .

Define  $c$  as follows:  $c(u_i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ;  $c(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j$  is even;  $c(v_j^i) = 1$  if  $i, j$  odd;  $c(v_j^i) = 2$  if  $i$  is even,  $j$  is odd; then  $\mathcal{S}(u_i) = 1$  if  $i$  is odd;  $\mathcal{S}(u_i) = 2$  if  $i$  is even;  $\mathcal{S}(v_j^i) = 1$  if  $i, j$  even;  $\mathcal{S}(v_j^i) = 2$  if  $i$  is odd,  $j$  is even;  $\mathcal{S}(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j$  is odd.

**Case 3.**  $n \equiv 0 \pmod{6}$ .

Define  $c$  as follows:

$c(u_i) = 0$  if  $i$  is odd;  $c(u_i) = 2$  if  $i$  is even;  $c(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j$  is even;  $c(v_j^i) = 1$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j$  is odd; then  $\mathcal{S}(u_i) = 1$  if  $i$  is odd;  $\mathcal{S}(u_i) = 0$  if  $i$  is even;  $\mathcal{S}(v_j^i) = 1$  if  $i, j$  even;  $\mathcal{S}(v_j^i) = 2$  if  $i$  is odd,  $j$  is even;  $\mathcal{S}(v_j^i) = 2$  if  $i$  is even,  $j$  is odd;  $\mathcal{S}(v_j^i) = 0$  if  $i, j$  odd. Clearly,  $\chi(C_m \circ C_n) = 3$ . Hence,  $Mc(C_m \circ C_n) = 3$ . This completes the proof.  $\square$

**Theorem 2.2.** For  $m$  even and  $n$  odd,  $m \geq 4$ ,  $n \geq 3$ ,  $Mc(C_m \circ C_n) = 4$ .

*Proof.* Let  $c : V(C_m \circ C_n) \rightarrow \mathbb{Z}_4$ .

**Case 1.**  $n \equiv 1 \pmod{8}$ .

Define  $c$  as follows:  $c(u_i) = 0$  if  $i$  is even;  $c(u_i) = 1$  if  $i$  is odd;  $c(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \equiv 0, 2, 3 \pmod{4}$ ;  $c(v_n^i) = c(v_{n-4}^i) = 1$  if  $i$  is odd;  $c(v_n^i) = 1$  if  $i$  is even;  $c(v_j^i) = 2$  if  $i$  is odd,  $j \equiv 1 \pmod{4}$ ;  $j \notin \{n, n-4\}$ ;  $c(v_j^i) = 2$  if  $i$  is even,  $j \equiv 1 \pmod{4}$ ,  $j \neq n$ ; then  $\mathcal{S}(u_i) = 0$  if  $i$  is odd;  $\mathcal{S}(u_i) = 3$  if  $i$  is even;  $\mathcal{S}(v_j^i) = 0$  if  $i$  is even,  $j \in \{3, 5, 7, \dots, n-2\}$ ;  $\mathcal{S}(v_j^i) = 1$  if  $i$  is odd,  $j \in \{3, 5, 7, \dots, n-2\}$ ;  $\mathcal{S}(v_j^i) = 1$  if  $i$  is even,  $j \in \{1, n-1\}$ ;  $\mathcal{S}(v_j^i) = 2$  if  $i$  is odd,  $j \in \{1, n-1, n-3, n-5\}$ ;  $\mathcal{S}(v_j^i) = 2$  if  $i$  is even,  $j \in \{2, 4, 6, \dots, n-3, n\}$ ;  $\mathcal{S}(v_j^i) = 3$  if  $i$  is odd,  $j \in \{2, 4, 6, \dots, n-7, n\}$ .

**Case 2.**  $n \equiv 3 \pmod{8}$  and  $n \neq 3$ .

Define  $c$  as follows:  $c(u_i) = 0$  if  $i$  is even;  $c(u_i) = 1$  if  $i$  is odd;  $c(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \equiv 0, 2, 3 \pmod{4}$ ;  $c(v_{n-2}^i) = c(v_{n-6}^i) = 1$  if  $i$  is odd;  $c(v_{n-2}^i) = 1$  if  $i$  is even;  $c(v_j^i) = 2$  if  $i$  is odd,  $j \equiv 1 \pmod{4}$ ;  $j \notin \{n-2, n-6\}$ ;  $c(v_j^i) = 2$  if  $i$  is even,  $j \equiv 1 \pmod{4}$ ,  $j \neq n-2$ ; then  $\mathcal{S}(u_i) = 0$  if  $i$  is odd;  $\mathcal{S}(u_i) = 3$  if  $i$  is even;  $\mathcal{S}(v_j^i) = 0$  if  $i$  is even,  $j \in \{1, 3, 5, \dots, n-2\}$ ;  $\mathcal{S}(v_j^i) = 1$  if  $i$  is odd,  $j \in \{1, 3, 5, \dots, n-2\}$ ;  $\mathcal{S}(v_j^i) = 1$  if  $i$  is even,  $j \in \{n-1, n-3\}$ ;  $\mathcal{S}(v_j^i) = 2$  if  $i$  is odd,  $j \in \{n-1, n-3, n-5, n-7\}$ ;  $\mathcal{S}(v_j^i) = 2$  if  $i$  is even,  $j \in \{2, 4, 6, \dots, n-5, n\}$ ;

$\mathcal{S}(v_j^i) = 3$  if  $i$  is odd,  $j \in \{2, 4, 6, \dots, n-9, n\}$ .

**Case 3.**  $n \equiv 5 \pmod{8}$ .

Define  $c$  as follows:  $c(u_i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ;  $c(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \equiv 0, 2, 3 \pmod{4}$ ;  $c(v_n^i) = 1$  if  $i$  is even;  $c(v_n^i) = 3$  if  $i$  is odd;  $c(v_j^i) = 2$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \equiv 1 \pmod{4}$ ,  $j \neq n$ ; then  $\mathcal{S}(u_i) = 1$  if  $i$  is odd;  $\mathcal{S}(u_i) = 3$  if  $i$  is even;  $\mathcal{S}(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \in \{3, 5, 7, \dots, n-2\}$ ;  $\mathcal{S}(v_j^i) = 1$  if  $i$  is even,  $j \in \{1, n-1\}$ ;  $\mathcal{S}(v_j^i) = 3$  if  $i$  is odd,  $j \in \{1, n-1\}$ ;  $\mathcal{S}(v_j^i) = 2$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \in \{2, 4, 6, \dots, n-3, n\}$ .

**Case 4.**  $n \equiv 7 \pmod{8}$ .

Define  $c$  as follows:  $c(u_i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ;  $c(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \equiv 0, 2, 3 \pmod{4}$ ;  $c(v_{n-2}^i) = 1$  if  $i$  is even;  $c(v_{n-2}^i) = 3$  if  $i$  is odd;  $c(v_j^i) = 2$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \equiv 1 \pmod{4}$ ,  $j \neq n-2$ ; then  $\mathcal{S}(u_i) = 1$  if  $i$  is odd;  $\mathcal{S}(u_i) = 3$  if  $i$  is even;  $\mathcal{S}(v_j^i) = 0$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \in \{1, 3, 5, \dots, n-2\}$ ;  $\mathcal{S}(v_j^i) = 1$  if  $i$  is even,  $j \in \{n-1, n-3\}$ ;  $\mathcal{S}(v_j^i) = 3$  if  $i$  is odd,  $j \in \{n-1, n-3\}$ ;  $\mathcal{S}(v_j^i) = 2$  if  $i \in \{1, 2, 3, \dots, m\}$ ,  $j \in \{2, 4, 6, \dots, n-5, n\}$ .

**Case 5.**  $n = 3$ .

**Subcase 5.1.**  $m \equiv 0 \pmod{4}$ .

Define  $c$  as follows:  $c(u_i) = 0$  if  $i \equiv 0, 2, 3 \pmod{4}$ ;  $c(u_i) = 1$  if  $i \equiv 1 \pmod{4}$ ;  $c(v_1^i) = 0$  if  $i \equiv 1 \pmod{4}$ ;  $c(v_2^i) = 2$  if  $i \equiv 1 \pmod{4}$ ;  $c(v_3^i) = 3$  if  $i \equiv 1 \pmod{4}$ ;  $c(v_1^i) = 0$  if  $i$  is even;  $c(v_2^i) = 1$  if  $i$  is even;  $c(v_3^i) = 2$  if  $i$  is even;  $c(v_1^i) = 1$  if  $i \equiv 3 \pmod{4}$ ;  $c(v_2^i) = 2$  if  $i \equiv 3 \pmod{4}$ ;  $c(v_3^i) = 3$  if  $i \equiv 3 \pmod{4}$ ; then  $\mathcal{S}(u_i) = 1$  if  $i \equiv 1 \pmod{4}$ ;  $\mathcal{S}(u_i) = 0$  if  $i \equiv 0, 2 \pmod{4}$ ;  $\mathcal{S}(u_i) = 2$  if  $i \equiv 3 \pmod{4}$ ;  $\mathcal{S}(v_1^i) = 2$  if  $i \equiv 1 \pmod{4}$ ;  $\mathcal{S}(v_2^i) = 0$  if  $i \equiv 1 \pmod{4}$ ;  $\mathcal{S}(v_3^i) = 3$  if  $i \equiv 1 \pmod{4}$ ;  $\mathcal{S}(v_1^i) = 1$  if  $i \equiv 3 \pmod{4}$ ;  $\mathcal{S}(v_2^i) = 0$  if  $i \equiv 3 \pmod{4}$ ;  $\mathcal{S}(v_3^i) = 3$  if  $i \equiv 3 \pmod{4}$ ;  $\mathcal{S}(v_1^i) = 3$  if  $i$  is even;  $\mathcal{S}(v_2^i) = 2$  if  $i$  is even;  $\mathcal{S}(v_3^i) = 1$  if  $i$  is even.

**Subcase 5.2.**  $m \equiv 2 \pmod{4}$ .

Define  $c$  as follows:  $c(u_i) = 0$  if  $i \equiv 0, 2, 3 \pmod{4}$ ;  $c(u_i) = 1$  if  $i \equiv 1 \pmod{4}$ ;  $c(v_1^i) = 0$  if  $i \equiv 1 \pmod{4}$ ;  $c(v_2^i) = 2$  if  $i \equiv 1 \pmod{4}$ ;  $c(v_3^i) = 3$  if  $i \equiv 1 \pmod{4}$ ;  $c(v_1^i) = 0$  if  $i \in \{2, 4, 6, \dots, m-2\}$ ;  $c(v_2^i) = 1$  if  $i \in \{2, 4, 6, \dots, m-2\}$ ;  $c(v_3^i) = 2$  if  $i \in \{2, 4, 6, \dots, m-2\}$ ;  $c(v_1^m) = 0$ ;  $c(v_2^m) = 1$ ;  $c(v_3^m) = 3$ ;  $c(v_1^i) = 1$  if  $i \equiv 3 \pmod{4}$ ;  $c(v_2^i) = 2$  if  $i \equiv 3 \pmod{4}$ ;  $c(v_3^i) = 3$  if  $i \equiv 3 \pmod{4}$ ; then  $\mathcal{S}(u_i) = 1$  if  $i \equiv 1 \pmod{4}$ ;  $\mathcal{S}(u_i) = 0$  if  $i \in \{2, 4, 6, \dots, m-2\}$ ;  $\mathcal{S}(u_i) = 2$  if  $i \equiv 3 \pmod{4}$ ;  $\mathcal{S}(u_m) = 2$ ;  $\mathcal{S}(v_1^i) = 2$  if  $i \equiv 1 \pmod{4}$ ;  $\mathcal{S}(v_2^i) = 0$  if  $i \equiv 1 \pmod{4}$ ;  $\mathcal{S}(v_3^i) = 3$  if  $i \equiv 1 \pmod{4}$ ;  $\mathcal{S}(v_1^i) = 1$  if  $i \equiv 3 \pmod{4}$ ;  $\mathcal{S}(v_2^i) = 0$  if  $i \equiv 3 \pmod{4}$ ;  $\mathcal{S}(v_3^i) = 3$  if  $i \equiv 3 \pmod{4}$ ;  $\mathcal{S}(v_1^i) = 3$  if  $i \in \{2, 4, 6, \dots, m-2\}$ ;  $\mathcal{S}(v_2^i) = 2$  if  $i \in \{2, 4, 6, \dots, m-2\}$ ;

$\mathcal{S}(v_3^i) = 1$  if  $i \in \{2, 4, 6, \dots, m-2\}$ ;  $\mathcal{S}(v_1^m) = 0$ ;  $\mathcal{S}(v_2^m) = 3$ ;  $\mathcal{S}(v_3^m) = 1$ . Clearly,  $Mc(C_m \circ C_n) \geq \chi(C_m \circ C_n) = 4$ . Hence,  $Mc(C_m \circ C_n) = 4$ . This completes the proof.  $\square$

### 3. CONCLUSION

For some graphs  $\mathcal{G}$  and  $\mathcal{H}$  considered in this paper, we have seen that  $Mc(\mathcal{G} \circ \mathcal{H}) = \chi(\mathcal{G} \circ \mathcal{H})$ . Except the case: For  $m \geq 1, n \geq 1$ ,  $Mc(C_{2m+1} \circ C_{2n+1})$ .

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